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# The DPRM-Theorem in Isabelle



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# The problem



*Eine diophantische Gleichung mit irgendwelchen Unbekannten und mit ganzen rationalen Zahlkoeffizienten sei vorgelegt: man soll ein Verfahren angeben, nach welchem sich mittels einer endlichen Anzahl von rationalen Operationen entscheiden lässt, ob die Gleichung in ganzen Zahlen lösbar ist.*

**DEF** A diophantine equation is a polynomial equation with integer coefficients

Examples:

$$5x - 10 = 0$$

$$x_1^2 - 4x_2 = x_3$$

$$x^3 + y^3 = z^3$$

# The problem



Is there an algorithm to determine, if a given diophantine equation has a solution in the integers?

**DEF** A diophantine equation is a polynomial equation with integer coefficients

Examples:

$$5x - 10 = 0$$

$$x_1^2 - 4x_2 = x_3$$

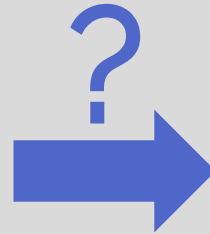
$$x^3 + y^3 = z^3$$

# Undecidability of Hilbert's problem



1900

diophantine

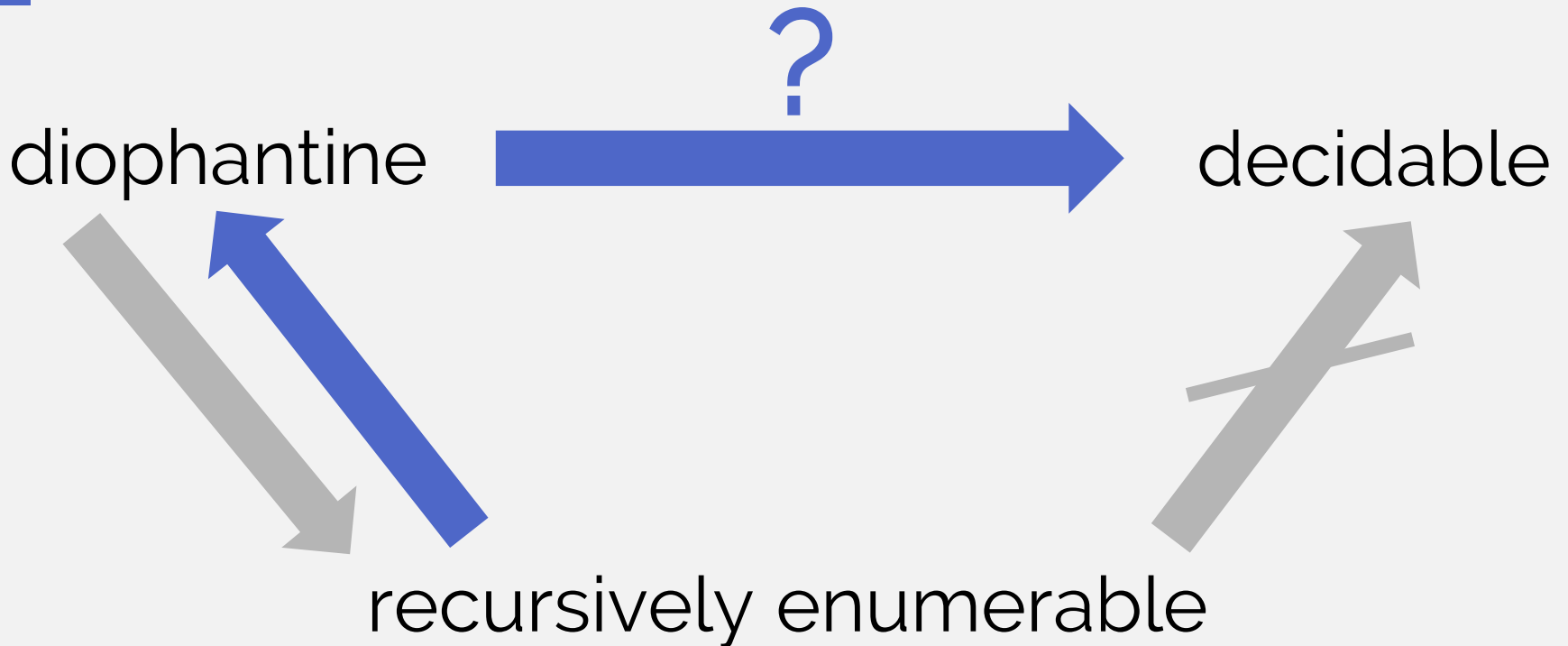


decidable

1950

Julia Robinson and Martin Davis conjecture undecidability

# Undecidability of Hilbert's problem



THM

DPRM Theorem.

Every recursively enumerable set is diophantine.

# Undecidability of Hilbert's problem



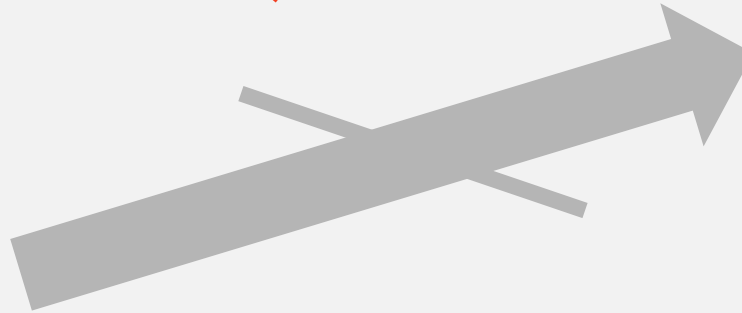
diophantine



recursively  
enumerable



decidable



THM

DPRM Theorem.

Every recursively enumerable set is diophantine.

# Undecidability of Hilbert's problem



diophantine



recursively  
enumerable



decidable

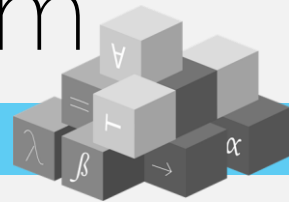


THM

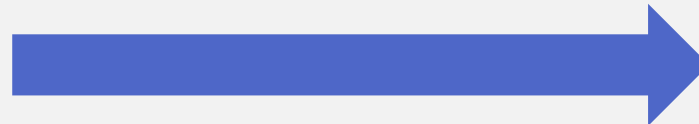
DPRM Theorem.

Every recursively enumerable set is diophantine.

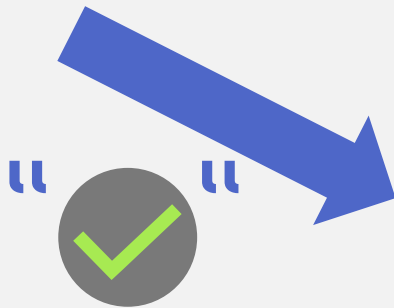
# Formalizing the DPRM Theorem



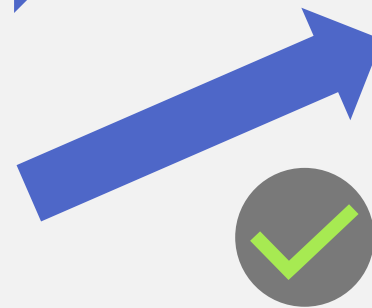
recursively  
enumerable



diophantine

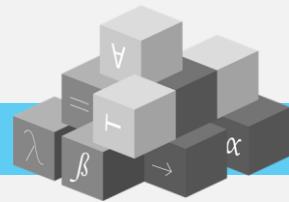


exponential  
diophantine

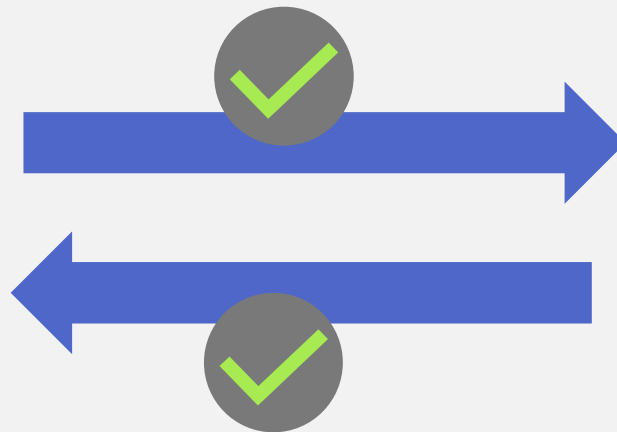




# Arithmetization



Register machine

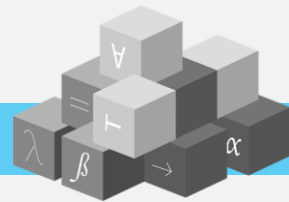


Arithmetization

up to positional  
notation lemmas

(exponential)  
diophantine

# Remaining work

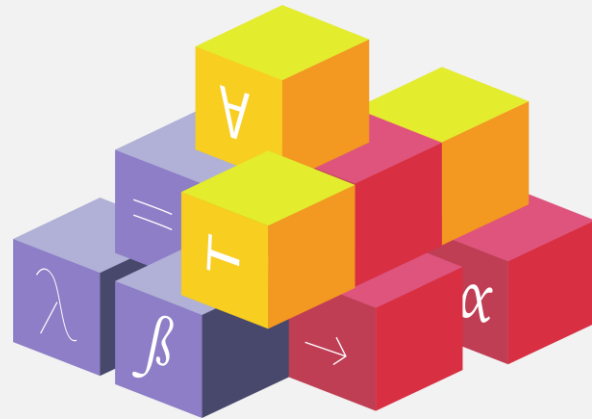


For conditions (4.6) and (4.7), it is very important that during summations there is no carry “across the boundary of a cell” in the protocol.

It can be shown by induction that for every  $t$ , among the numbers  $s_{1,t}, \dots, s_{m,t}$  one and only one is equal to 1, others being equal to 0. This is so for  $t = 0$  thanks to (4.24) and (4.25). If this is so for some  $t$ , then the same holds for  $t + 1$ . This is so because in the right-hand side of (4.7), all summands but 1 should be equal to 0, the remaining summand being equal to 1; hence, no carry at all occurs in summations (4.24)–(4.25) and relation (4.7) holds.

Similarly, in (4.6), besides the first summand, there can be at most one other summand different from 0. By (4.15), the first summand is at most  $2^c - 1$ , and the other non-zero summand can be equal only to 1. This implies that no carry “across the boundary of a cell” ever occurs and, hence, relation (4.6) also holds.

Questions?



# Thank you for your attention!

And a lot of thanks to

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